

Quality Meshing of a Forest of Branching Structures

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joint work with

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Outline

- 1 Motivation and Problem Statement
- 2 Prior Work and Background
- 3 Approach
- 4 Initial Results

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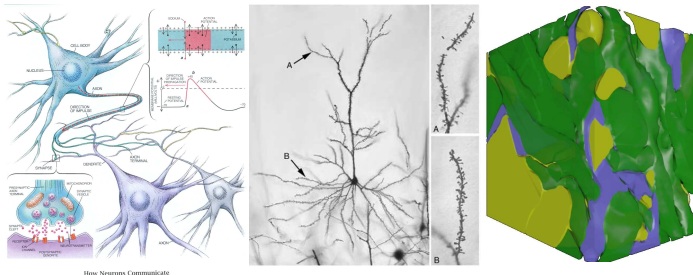
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Motivation: Neuronal Modeling



Neurons *in vivo* are packed very densely and have many small geometric features known to affect voltage decay.

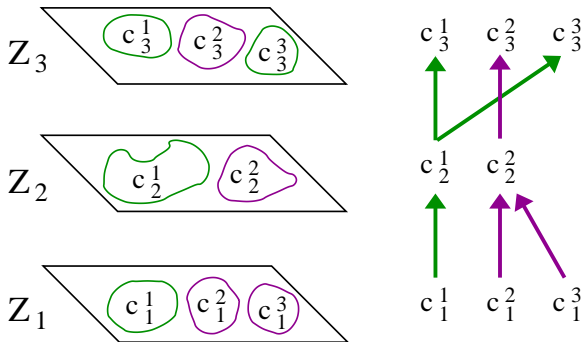


Neuron length $\sim 100 \mu\text{m}$; Neuropil dataset $\sim 2 (\mu\text{m})^3$; in-plane resolution $\sim 5\text{-}10 \text{ nm}$

Formal Problem Statement

Input (for a K component forest):

- 1 M horizontal planes Z_1, \dots, Z_M . (Z_m given by $z = z_m$)
- 2 K functions $g_m^1, \dots, g_m^K : \mathbb{R}^2 \rightarrow \mathbb{R}$ such that $\bigcup_{k=1}^K \{g_m^k = 0\}$ is a 1-manifold.
- 3 Contours $\{c_m^k\}$ of the set $\{g_m^k = 0\}$.
- 4 An acyclic directed graph G with vertices $\{c_m^k\}$, indicating connectivity.



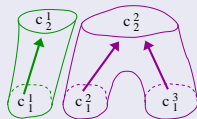
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Output K functions $h_1, \dots, h_K : \mathbb{R}^3 \rightarrow \mathbb{R}$ such that:

- 1 Each h_k restricts to g_m^k on Z_m , i.e.
 $h_k(x, y, z_m) \equiv g_m^k(x, y)$.
- 2 Each surface $h_k(x, y, z) = 0$ is a compact, connected, smooth 2-manifold with local connectivity corresponding to the graph G .
- 3 The K component surface $\prod_{k=1}^K h_k(x, y, z) = 0$ is a 2-manifold, i.e. the component surfaces do not intersect.



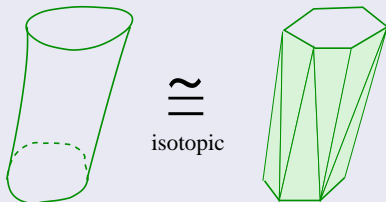
Simplifying Assumptions

Assumptions for Reduction to Meshing:

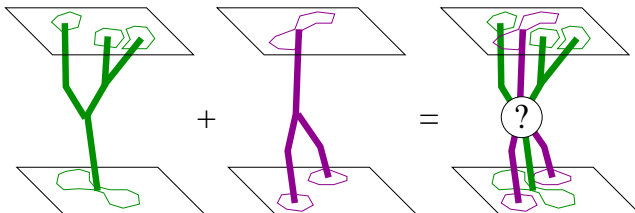
- 1 Contours are simple polygons and can be refined if necessary.



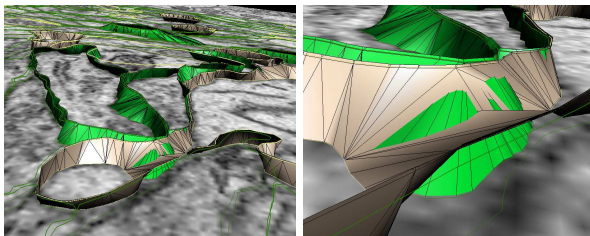
- 2 A mesh of the polygonal contours satisfying the output properties is isotopic to a smooth solution.



The Multi-Component Difficulty



Independent solutions to the reconstruction problem for each component may produce topological or geometrical inaccuracies when aggregated.



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Selected Prior Work on Single Component Problem

FUCHS, KEDEM AND USELTON *Optimal surface reconstruction from planar contours* Communications of the ACM 20:10, 1977.

- Seminal work in reconstruction from polygonal contours.

MEYERS, SKINNER AND SLOAN *Surfaces from contours* ACM Transactions on Graphics 11:3, 1992.

- Identified subproblems of correspondence, tiling, and branching.

BAREQUET AND SHARIR *Piecewise-linear interpolation between polygonal slices* Symposium on Computational Geometry 1994.

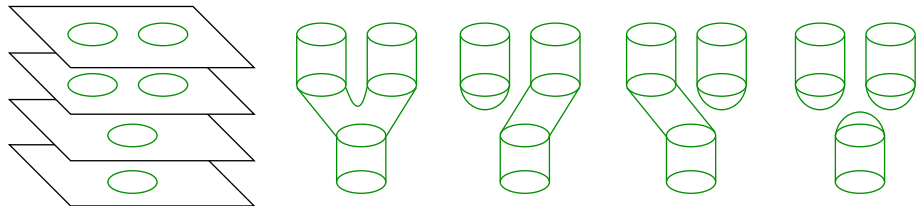
- Developed an algorithm for CT, MRI, and other medical applications.

BAJAJ, COYLE AND LIN *Arbitrary topology shape reconstruction from planar cross sections* Graphic Models and Image Processing 58:6, 1996.

- Expanded this algorithm by providing topological guarantees on the output.
- We use this method for our approach.

The Correspondence Problem

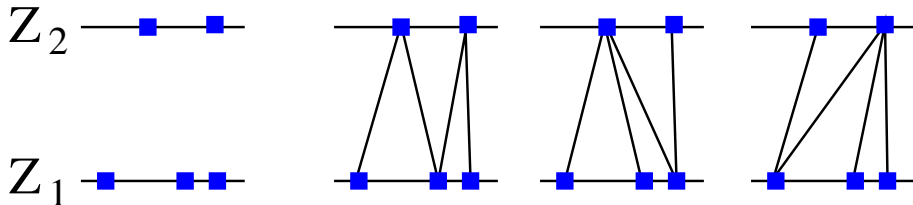
How should contours on adjacent slices connect?



Remark: This is resolved by the connectivity graph G in our case.

The Tiling Problem

How should corresponding contours be tiled?



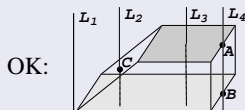
Definition: A **slice chord** is an edge connecting vertices on adjacent slices. A **tiling triangle** is formed by two slice chords and a contour edge.

What are the criteria of a “good” tiling?

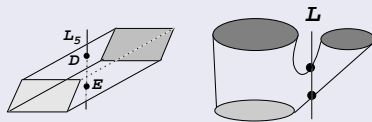
The Tiling Problem

Desired Tiling Criteria

- 1 The reconstructed surface forms a piecewise closed surface of polyhedra.
- 2 Any vertical line segment between two adjacent slices intersecting the reconstructed surface does so at exactly one point or along exactly one line segment.



not OK:

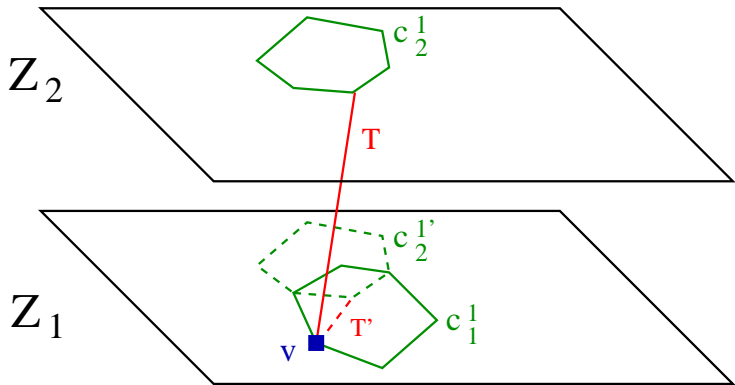


- 3 Resampling of the reconstructed surface on any slice reproduces the original contours.

Remark: The last criterion implies that aside from contour refinement, any edges or vertices added to the mesh must lie outside the Z_i planes.

The Tiling Problem Resolved

- Let v be a vertex in contour $c_1^1 \subset Z_1$ corresponding to $c_2^1 \subset Z_2$.
- Let T be a slice chord from v to c_2^1 .
- Let the "prime" notation denote vertical projection to Z_1 .

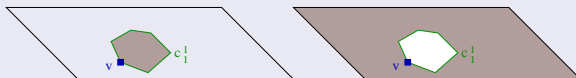


The Tiling Problem Resolved

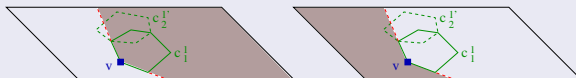
Theorem [Bajaj, Coyle, Lin 1996]:

If a tiling satisfies the three criteria, the following hold:

i) T' lies in exactly one of these regions:



ii) If $v \notin c_2^{1'}$ then T' lies in exactly one of these regions:



iii) If $v \in c_2^{1'}$ then T' lies in exactly one of these regions:



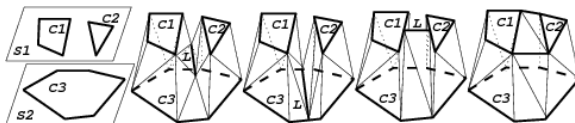
The Tiling Problem Resolved

Tiling Algorithm [*Bajaj, Coyle, Lin 1996*]:

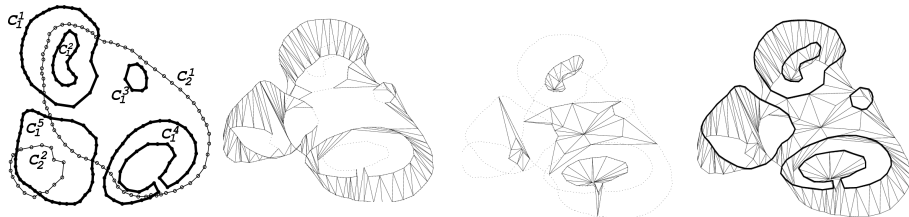
- 1 For each vertex $v \in \{c_1^j\}$, make a list of all the slice chords that could be formed to a vertex of $\{c_2^k\}$ (based on the resolution of the correspondence problem).
- 2 Select the shortest length chord from this list which satisfies the results of the Theorem.
- 3 If no chord from the list satisfies the theorem, tag the vertex as “untiled.”
- 4 Collect boundaries of untiled regions for subsequent meshing when resolving the branching problem.

The Branching Problem

How should tiling be done when a contour in Z_1 corresponds to more than one contour in Z_2 ?

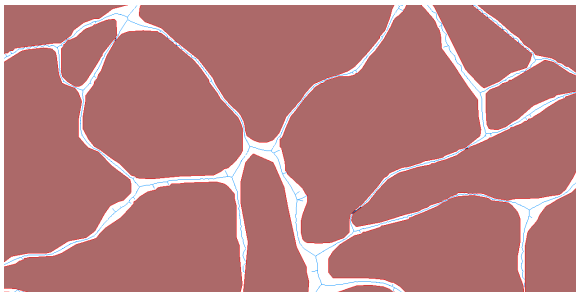


To ensure the criteria are satisfied, we add vertices to a plane half way between Z_1 and Z_2 and then mesh.



The Medial Axis

For our approach, we will use the medial axis of the region exterior to the contours on each slice.



Definition: The **medial axis** \mathfrak{M} of an open set $O \subset \mathbb{R}^n$ is the set of points $x \in O$ for which there are at least two closest points to x on the complement O^c .

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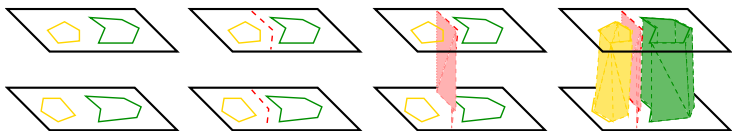
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Algorithm

Idea: Compute the medial axis on each slice, scaffold it to a three dimensional partitioning, and use this to detect and remove intersections between components.

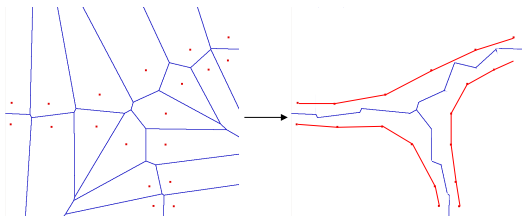


$\text{MULTISURFRECON}(\{Z_i\}, \{c_i^j\}, \{P_i\}, G)$ (P_i = vertices of the contours c_i^j)

- 1 Construct and partition medial axes on each slice.
- 2 Construct a medial surface.
- 3 Create single component surfaces.
- 4 Remove overlaps on intermediate planes.
- 5 Improve mesh quality.

For subsequent explanation, suppose G has only two components, yellow and green, so that each c_i^j is either a yellow contour (YC) or green contour (GC).

- 1 Construct and partition medial axes on each slice.
- Compute the Voronoi diagram $\text{Vor } P_i$ of the sample points P_i of the contours on slice i .
- Discard all Voronoi edges intersecting the contours.
- Discard all remaining vertices and edges interior to the contours.
- The remaining vertices and edges approximate the medial axis \mathfrak{M} of extracellular space.
- Let $E_{GY} = \{\text{edges of } \mathfrak{M} \text{ defined by a point of YC and a point of GC}\}$

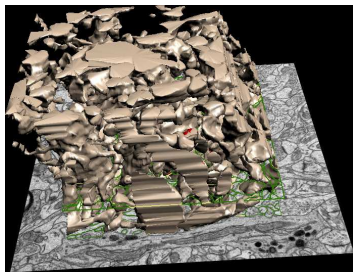
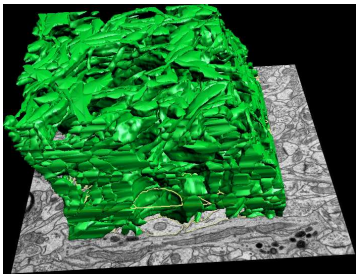


Algorithm

2 Construct a medial surface.

Dilate the edges in E_{GY} so that they form contours and color these contours red. Create a partial, approximate medial surface \mathcal{MS} by running the SINGLESURFRECON code on E_{GY} .

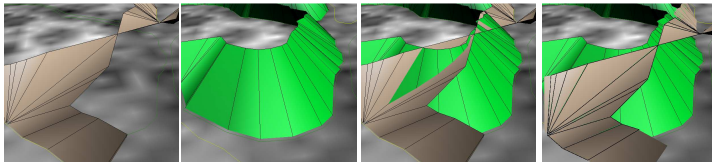
3 Create single component surfaces.



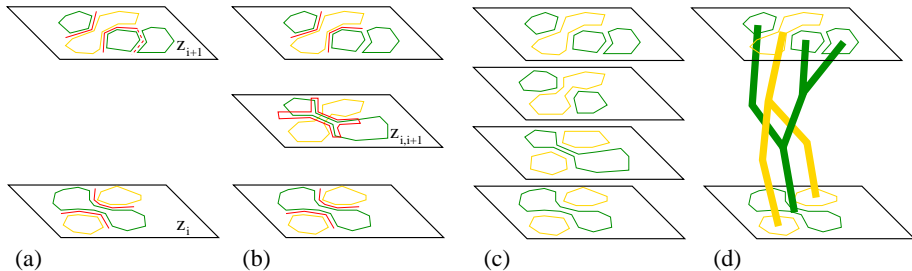
Algorithm

4 Remove overlaps on intermediate planes.

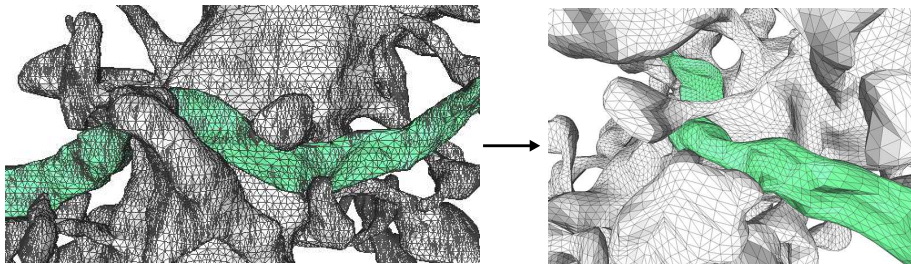
Simple overlap:



Exotic overlap:



5 Improve mesh quality.



- Decimation: QSlim software by Michael Garland.
- Shape Improvement: Geometric Flow library from LBIE Mesher, part of Volume Rover software by CVC lab.

Correctness Results

Lemma 1:

Suppose that for each i , there are no pairwise overlaps between the original contours on plane Z_i . Then the output of SINGLESURFRECON run on any subset of the original contours is not self-intersecting.

Proof: This is a consequence of the three tiling criteria; if the original contours do not overlap, SINGLESURFRECON cannot create self-intersections.

Lemma 2:

If there are no overlaps of any of the intermediate planes where branching occurs, then the output of MULTISURFRECON has no self-intersections in the entire volume. Conversely, if the output has no self-intersections in the volume, it does not have any overlap on any intermediate plane.

Proof: \Rightarrow : The output is a linear interpolation between consecutive planes. If the contours do not overlap, their interpolation does not intersect.

\Leftarrow : If the output is not self-intersecting, the components cannot intersect on any plane.

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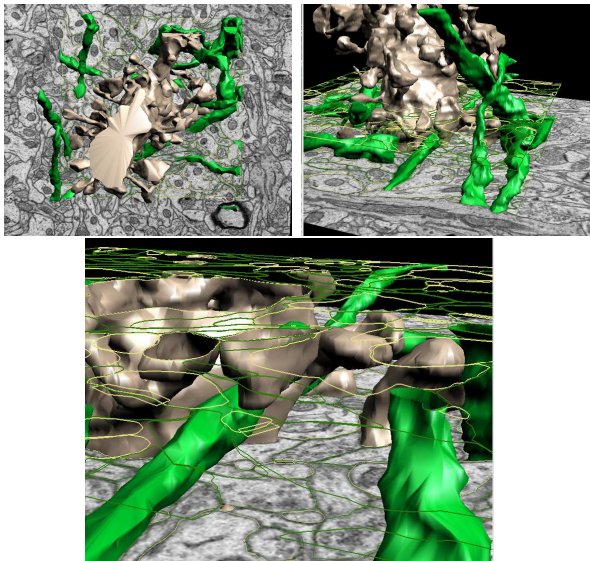
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Initial Results



Acknowledgements



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