

407
10/19

Problems: 12, 13.

LL

* Change Terms.

Def: Let f be a real function, and S a subset of its domain. We say that f is monotone

$$\begin{cases} \text{- increasing on } S \text{ iff } x < y \Rightarrow f(x) \leq f(y), \forall x, y \in S \\ \text{- decreasing on } S \text{ iff } x < y \Rightarrow f(x) > f(y), \forall x, y \in S. \end{cases}$$

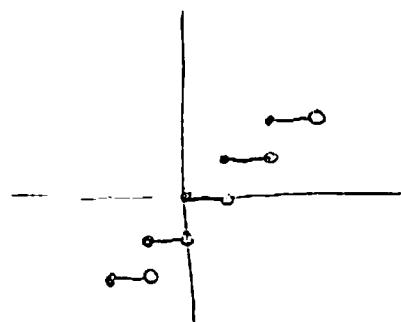
We say f is strictly $\begin{cases} \text{increasing} \\ \text{decreasing} \end{cases}$ on S iff

$$\forall x < y \Rightarrow \begin{cases} f(x) < f(y) \\ f(x) > f(y) \end{cases} \quad \forall x, y \in S.$$

and that f is strictly monotone if it is strictly inc or dec.

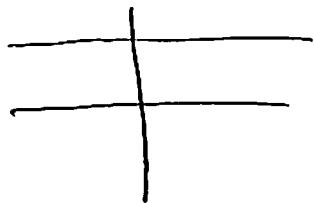
Example:

1) $f(x) = \lfloor x \rfloor := \left\{ \begin{array}{l} \text{the largest integer} \\ \leq x \end{array} \right\}$ is monotone increasing,
but not strictly increasing

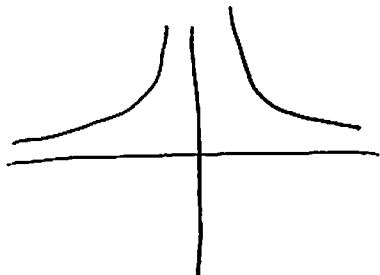


[2]

2) $f(x) = c$ (constant) is monotone increasing
and monotone decreasing,
 but neither increasing nor
 decreasing (not strictly
 monotone)



3) $f(x) = \frac{1}{x^2}$ is Strictly increasing on $\mathbb{R}_{<0}$
 Strictly decreasing on $\mathbb{R}_{>0}$
 "witch hat".



Thm: If f is strictly monotone on S , then f is H on S .

proof: Let $x, y \in S$ and suppose $x \neq y$. Then either $x < y$ or $y < x$, and wlog $x < y$. Since f is strictly monotone on S , this implies $f(x) < f(y)$ or $f(x) > f(y)$; in either case $f(x) \neq f(y)$ so f is H on S .

Cor: If f is strictly monotone on S , then $f|_S$ has an inverse.

Now suppose that f is a differentiable function on $S \subseteq \mathbb{R}$.

Thm: - If $f'(x) > 0 \quad \forall x \in S$ then f is strictly increasing on S

- If $f'(x) < 0 \quad \forall x \in S$ then f is strictly decreasing on S .

Pf. Suppose that $f(x) = f(y)$ for $x, y \in S$ with $x \neq y$ ($\log_{x \neq y}$)

Then by the mean value theorem, $\exists \alpha \in (x, y)$

$$\text{s.t. } f'(\alpha) = \frac{f(y) - f(x)}{y - x} = 0, \text{ Thus, if } f'(x) \neq 0 \quad \forall x \in S$$

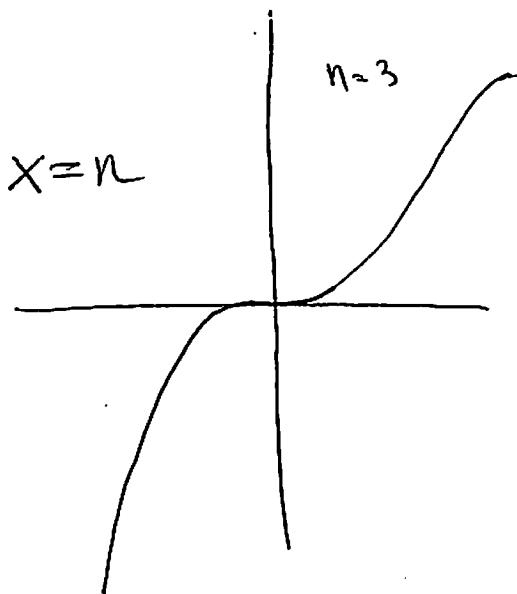
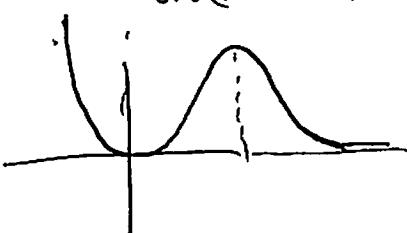
then f is strictly monotone on S .

Ex: $f(x) = x^n e^{-x}$ is not monotone $\forall n > 0$ on \mathbb{R} .

$$f'(x) = nx^{n-1}e^{-x} - x^n e^{-x}$$

$$= e^{-x} x^{n-1} (n - x) = 0 \quad \text{for } x = n$$

and $f' > 0$ if $0 < x < n$
 < 0 if $x > n$



L4

The derivative condition may also be interpreted as follows:

f differentiable on S , $x_0 \in S$, then

$$f(x) = f(x_0) + f'(x_0)(x-x_0) \quad \text{Provided } f'(x_0) \neq 0,$$

$$\text{So } \frac{(f(x) - f(x_0))}{f'(x_0)} + x_0 \cong x \quad \text{and}$$

$$f^{-1}(x) \cong \frac{x - f(x_0)}{f'(x_0)} + x_0$$

Thm: If f is strictly $\begin{cases} \text{incr.} \\ \text{decr} \end{cases}$ on S

then f^{-1} is strictly $\begin{cases} \text{incr} \\ \text{decr} \end{cases}$ on $f(S)$.

Pf: If f is strictly incr, then for $x, y \in f(S)$,

$$\begin{aligned} f^{-1}(x) < f^{-1}(y) &\Rightarrow f(f^{-1}(x)) < f(f^{-1}(y)) \\ &\Rightarrow x < y \end{aligned}$$

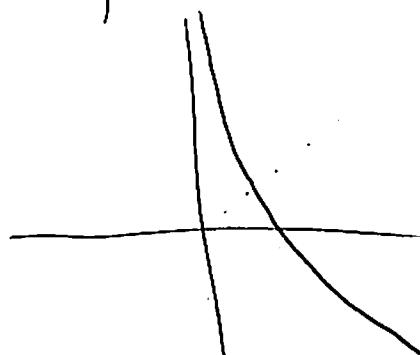
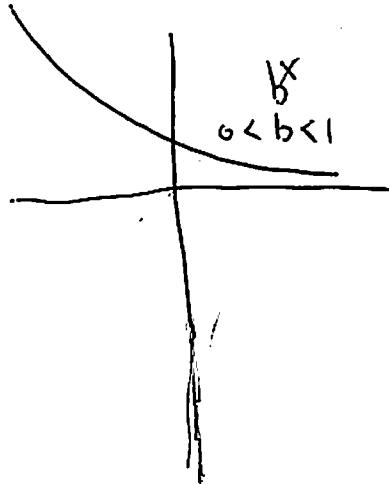
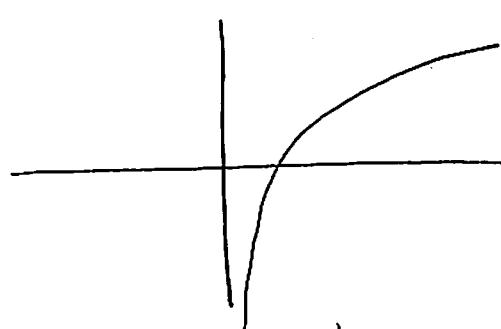
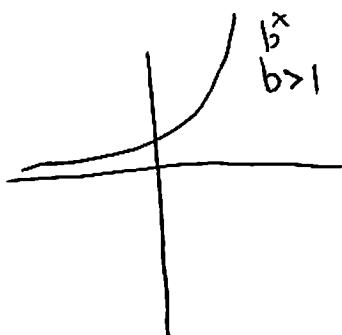
Hence, $x > y \Rightarrow f^{-1}(x) > f^{-1}(y)$

Since $x = y$ iff $f^{-1}(x) = f^{-1}(y)$, we conclude
 $x > y \Rightarrow f^{-1}(x) > f^{-1}(y)$. Similar pf for str. dec.

Ex: $b \neq 1$. Then $y = f(x) = b^x$
 is strictly $\begin{cases} \text{increasing} \\ \text{decreasing} \end{cases}$ if $\begin{cases} 1 < b \\ 0 < b < 1 \end{cases}$.

We have $f^{-1}(x) = \log_b(x)$. It is

strictly $\begin{cases} \text{increasing} \\ \text{decreasing} \end{cases}$ if $\begin{cases} 1 < b \\ 0 < b < 1 \end{cases}$



$$\log_b(x) = -\log_{\frac{1}{b}}(x)$$

Sequences:

$\{S_n\}_{n \in \mathbb{Z}_{\geq k}}$ a sequence.

$\{S_n\}$ is $\begin{cases} \text{increasing} \\ \text{decreasing} \end{cases}$ iff $\begin{cases} S_n < S_{n+1} \\ S_n > S_{n+1} \end{cases} \quad \forall n \in \mathbb{Z}_{\geq k}$.

Ex: Let s_n be defined by $s_0 = 1$,

$$s_{n+1} = \frac{1}{8}(3s_n + 6).$$

We compute:

n	1	2	3	4	5
s_n	1	$\frac{9}{8}$	$\frac{75}{64}$	\dots	

Claim s_n is ~~strictly~~ increasing.

Pf. ~~Induction~~. $s_{n+1} = \frac{1}{8}(3s_n + 6) > s_n$

$$\Leftrightarrow 3s_n + 6 > 8s_n$$

$$\Leftrightarrow 6 > 5s_n$$

$$\Leftrightarrow s_n < \frac{6}{5} \quad \forall n.$$

We prove this by induction:

base case: $s_0 = 1 < \frac{6}{5}$ ✓

If $s_k < \frac{6}{5}$ then $s_{k+1} = \frac{1}{8}(3s_k + 6)$

$$\begin{aligned} &< \frac{1}{8}(3 \cdot \frac{6}{5} + 6) \\ &= \frac{1}{8}(\frac{48}{5}) = \frac{6}{5}. \quad \blacksquare \end{aligned}$$

Ex: Take any $n \in \mathbb{Z}_{\geq 1}$. Define $\{s_i(n)\}$ by

$$s_1(n) = n, \quad s_{i+1}(n) = \begin{cases} 3s_i(n) + 1 & \text{if } s_i(n) \text{ is odd} \\ \frac{s_i(n)}{2} & \text{if } s_i(n) \text{ is even} \end{cases}$$